# Efficient Benchmarking with Provable Guarantees and more

#### Sami Abdul Sater

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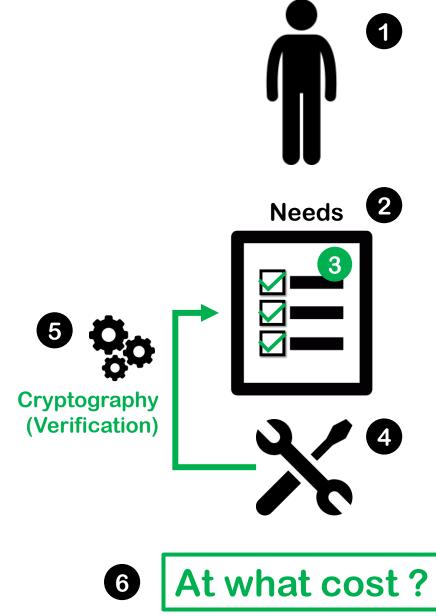






### Outline

- Motivation
  - 1. Benchmark... for who?
  - 2. Addressing which needs?
- Contributions
  - What guarantees are reached?
  - Benchmarking protocol
  - **Origins: Verification**
  - Consequence on assumptions
- Conclusion



At what cost?

### Motivation: benchmark... for whom?

Benchmark: An approach that enables an entity to compare options based on a metric that is relevant to its usage.

• Enlighten hardware constructors...







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Benchmark: An approach that enables an entity to compare options based on a metric that is relevant to its usage.

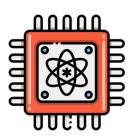
Enlighten hardware constructors...

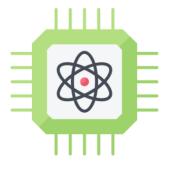
#### THIS WORK:

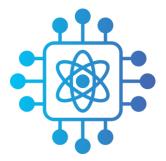
Enlighten hardware buyers.

Very generic usage!









# Addressing HW buyers needs

HW = hardware

#### Wish-list\*

- ☐ Reliable <u>QC capability</u>
- ☐ Abstraction on the HW











#### **Benchmark properties**

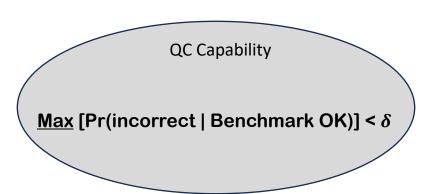
- Scalability: what happens when the target computation size gets bigger?
- Efficiency: what happens when we want a better guarantee (smaller  $\delta$ )?
- Minimal assumptions

(technology-independent, ...)

#### Output of a benchmark: Meaningful\* guarantee

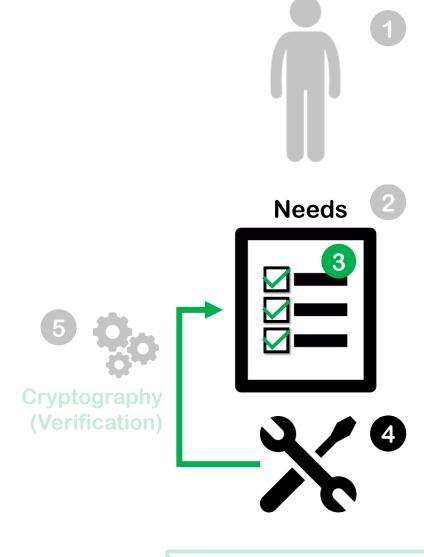


 $\delta$  –soundness Worst—case guarantee



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### Technical statements



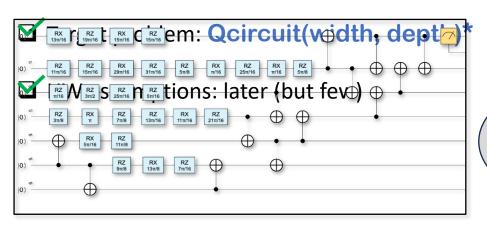




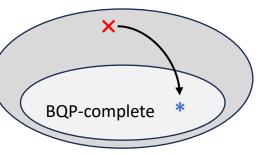




#### **Evaluating a Quantum Computer** (on useful problems)



BQP (quantum BPP)



#### **Benchmark properties**

**Efficiency** 

#### **Output of the benchmark:**

- Soundness:  $Pr(incorrect | Benchmark OK) < \delta$
- **Proved** by cryptography
- Worst-case guarantee: for all instances\*, in particular the worst one

efficient

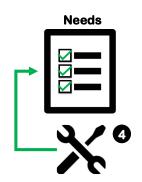
# "test" rounds:

 $\sim log(1/\delta)$ 

Size of "test": size of the comp.

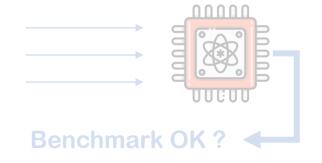
scalable

→ Meaningful and relevant guarantee for HW buyers



1. Define guarantee  $\delta$  — Number (#) of **test rounds** + size (depth, width) of comp.

- $\# \sim \log(1/\delta)$  $\delta \sim \exp(-\#)$
- 2. Test phase. Run **tests**, analyze outcomes, decide if **Benchmark OK** or not



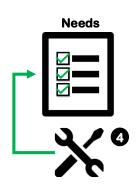
Target comp.

Test.

Check!

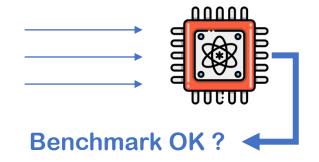
(different inputs, same comp, deterministic outcomes)

3. Guarantee. For any instance of QCircuit of the same width and depth,



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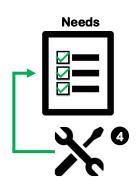
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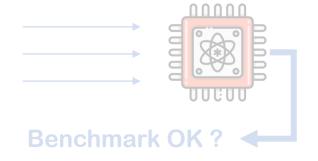
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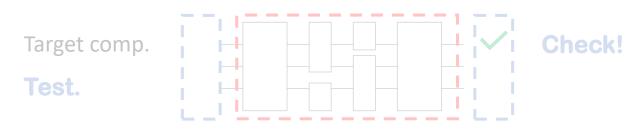
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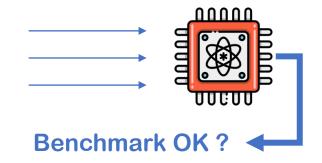


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Target comp.

Test.

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(different inputs, same comp, deterministic outcomes)

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#### **BENCHMARK PROPERTIES**

#### Meaningful guarantee

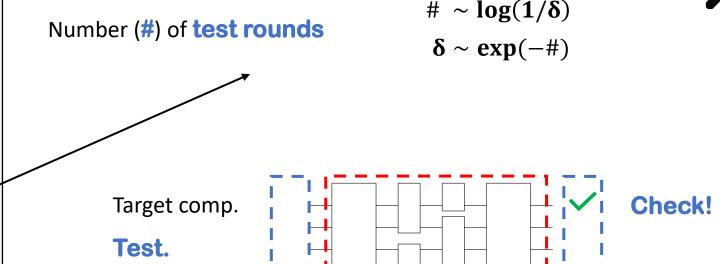
- Bound on the failure probability
- For any computation of the class
- Proved bound

Efficiency: more precision?

More tests, but logarithmic

Scalability: bigger computation?

Only impacts tests' size (!)

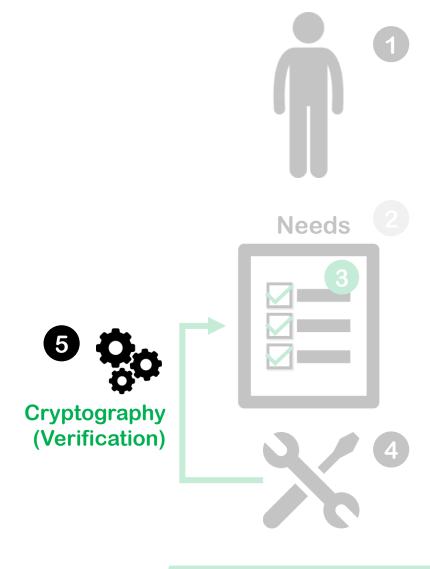


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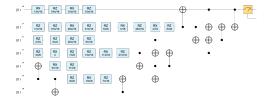
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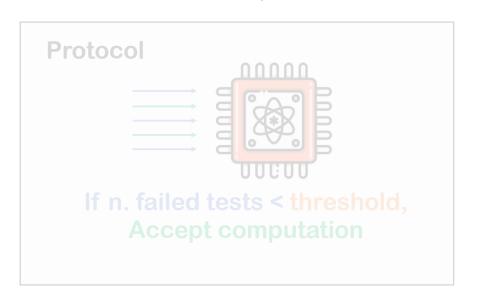


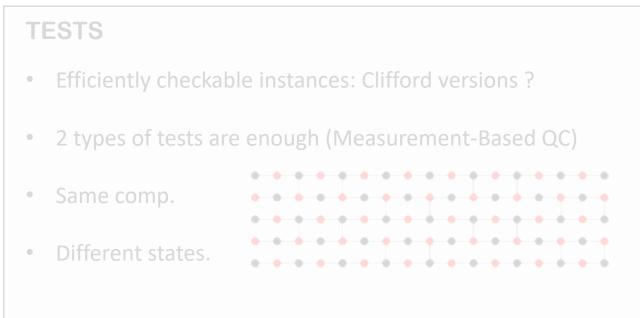
1. Define



# test rounds, computation rounds, threshold

2. Interleave test and computation rounds.





Guarantee. For any instance of QCircuit of the same width and depth,

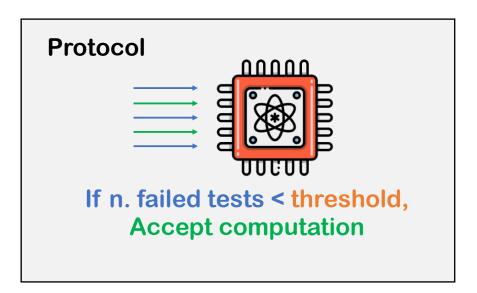
Pr(Failed computation | n. failed tests < threshold) <  $\delta$ 



1. Define

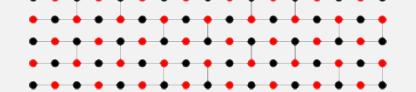
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#### **TESTS**

- Efficiently checkable instances: Clifford versions?
- 2 types of tests are enough (Measurement-Based QC)
- Same comp.
- Different states.



3. Guarantee. For any instance of QCircuit of the same width and depth,

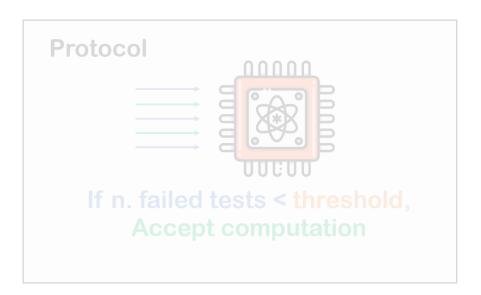
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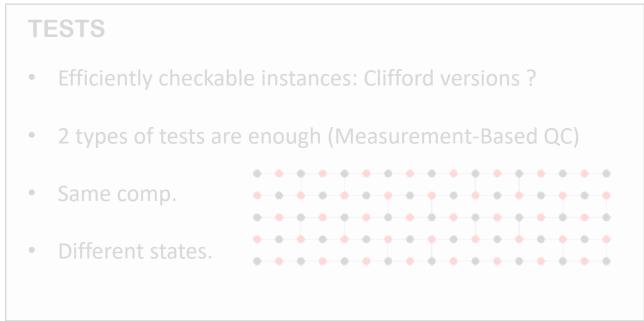


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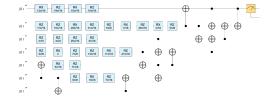


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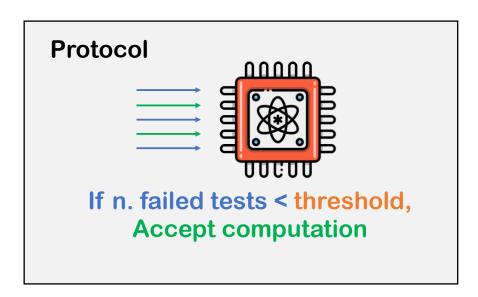


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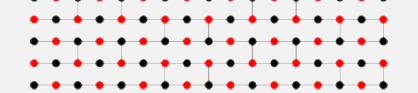
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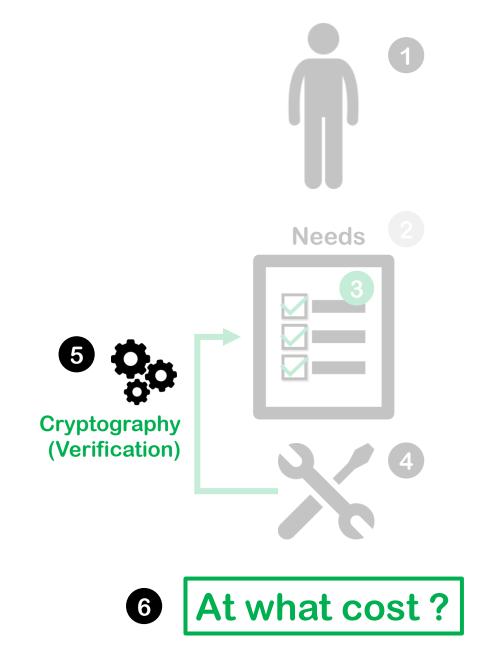
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Pr(Failed computation | n. failed tests < threshold) <  $\delta$ 

 $\delta \in \text{negl}(\#, \#)$ 

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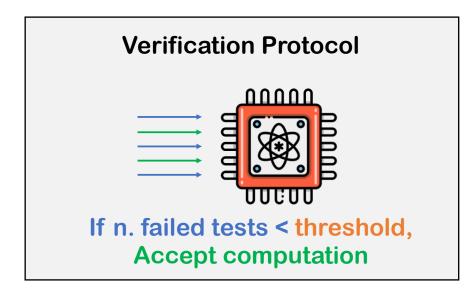


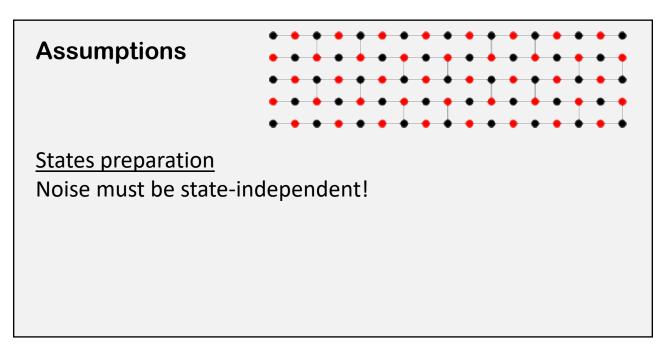
### Benchmark via Verification: at what cost?

1. Define

# test rounds, computation rounds, threshold

2. Interleave test and computation rounds.





3. Guarantee. For any instance of QCircuit of the same width and depth,

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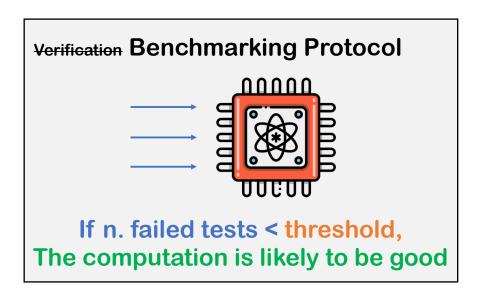
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### Benchmark via Verification: at what cost?

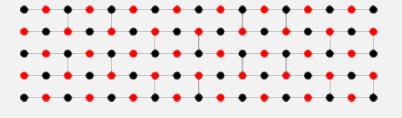
1. Define

# test rounds, computation rounds, threshold

2. Interleave test and computation rounds.



**Assumptions** 



**States preparation** 

Noise must be state-independent!

Repeatability

HW behaves the same through rounds

3. Guarantee.

For any instance of QCircuit of the same width and depth, when we do the computation,

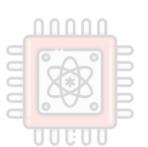
Pr(Failed computation | n. failed tests < threshold) <  $\delta$  (Benchmark OK)

 $\delta \in \text{negl}(\#,\#)$ 

### Conclusion: what does it tell us?

- 1. A protocol evaluating a QC's performance through a metric characterizing its reliability on a set of useful computations, relevant for HW Buyers
- 2. Efficient: repetition is the only overhead
- 3. Scalable: the computation size only affects the tests size
- 4. The guarantee is valid under assumptions.
  - Repeatability
  - State-independent noise
  - Target for HW constructors!
- And more!
  - Simpler, circuit-model protocol
  - With assumption on the noise model, possibility to do noise characterization







### Thank you for your attention!



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### Backslides

